

MAT2440, Classwork7, Spring2025

ID: _____

Name: _____

1. Use identities to prove " $\neg(p \rightarrow q) \equiv p \wedge \neg q$ "

$$\neg(p \rightarrow q) \equiv p \wedge (\neg q)$$

$$\neg(p \rightarrow q) \equiv \neg(\neg p \vee q)$$

first one in Group II

$$\equiv \neg(\neg p) \wedge \neg q$$

De Morgan's law

$$\equiv p \wedge (\neg q)$$

Double Negation

2. Use identities to prove " $(p \wedge q) \rightarrow (p \vee q)$ " is a tautology.

$$(p \wedge q) \rightarrow (p \vee q) \equiv \neg(p \wedge q) \vee (p \vee q)$$

$$\equiv (\neg p \vee \neg q) \vee (p \vee q)$$

De Morgan's

$$\equiv (\neg p \vee p) \vee (\neg q \vee q)$$

associative and
commutative

$$\equiv \top \vee \top$$

negation

$$\equiv \top$$

domination

3. Group III of the logically equivalences: Identities related to biconditional statements.

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p) \leftarrow \text{"definition" of "p \leftrightarrow q"}$$

$$p \leftrightarrow q \equiv \neg p \leftrightarrow \neg q$$

$$(3) p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$$

$$\neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$$

$$\text{Prove } "p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)"$$

$$\text{Proof: } p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

$$\equiv (\neg p \vee q) \wedge (\neg q \vee p)$$

Group II, the 1st one

$$\equiv ((\neg p \wedge \neg q) \vee (q \wedge \neg q)) \vee ((\neg p \wedge p) \vee (q \wedge p))$$

negation law

$$\equiv ((\neg p \wedge \neg q) \vee F) \vee (F \vee (q \wedge p))$$

Identity law

$$\equiv (\neg p \wedge \neg q) \vee (q \wedge p)$$

$$\equiv (p \wedge q) \vee (\neg p \wedge \neg q)$$

$$\text{prove } \neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$$

proof: Group III (3)

$$\neg(p \leftrightarrow q) \equiv \neg((p \wedge q) \vee (\neg p \wedge \neg q))$$

De Morgan's

$$\equiv (\neg(p \wedge q)) \wedge \neg(\neg p \wedge \neg q)$$

$$\equiv (\neg p \rightarrow \neg q) \wedge (p \vee q)$$

De Morgan's

$\neg(p \rightarrow \neg q) \equiv p \wedge \neg q$

and negation

in Group II

$$\equiv (\neg p \rightarrow \neg q) \wedge (p \vee \neg(\neg q))$$

$$\equiv (\neg p \rightarrow \neg q) \wedge (\neg q \rightarrow p)$$

(Group II: $p \rightarrow q \equiv (\neg p) \vee q$)

$$\equiv p \leftrightarrow \neg q \quad (\text{Group III first one})$$

4. Predicate logic and Propositional function:

The predicate logic allows variables in propositions and enables us to reason and explore relationships between objects. A Propositional function is a statement with variables and has been used on predicate logic. Once the values have been assigned to the variables, the propositional function becomes a proposition, and has truth value,

5. Let $P(x)$ denote the statement " $x > 3$ ". What are the truth values of $P(2)$ and $P(4)$.

$P(x)$ denotes " $\underbrace{x}_{\substack{\uparrow \\ \text{Variable} \\ (\text{subject})}} \text{ is greater than } \underbrace{3}_{\substack{\uparrow \\ \text{predicate} \\ (\text{property})}}$ "

$P(2)$ is " $2 > 3$ " false
 $P(4)$ is " $4 > 3$ " True

6. Let $Q(x, y)$ denote the statement " $x = y + 3$ ". What are the truth values of the propositions

$Q(1, 2)$ and $Q(3, 0)$.

$Q(1, 2)$ means " $1 = 2 + 3$ " which is false.

$Q(3, 0)$ means " $3 = 0 + 3$ " which is true.

7. Given a computer programming "If $\underbrace{x}_{\substack{\uparrow \\ \text{variable}}} > 0$, then $x := x + 1$ ". Using the terminology of propositional function to explain it.

this is $P(x)$

If $P(x)$ is true, then x is increased by 1.

If $P(x)$ is false, then x is not changed.