

# MAT2440, Classwork43, Spring2025

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## 1. Affine Cipher.

The shift ciphers can be generalized further to slightly enhance security by using a function of the form

$$f(p) = \underline{ap + b} \pmod{26},$$

where  $a$  and  $b$  are integers and  $\underline{\gcd(a, 26) = 1}$  to ensure the existence of the decryption function. Such a cipher is called affine cipher.

## 2. Encrypt the plaintext "PARK" when function $f(p) = 7p + 3 \pmod{26}$ is used.

|    |   |    |    |   |
|----|---|----|----|---|
| P  | A | R  | K  | $f(15) = 7 \cdot 15 + 3 \pmod{26} = 108 \pmod{26} = 4$  |
| 15 | 0 | 17 | 10 | $f(0) = 7 \cdot 0 + 3 \pmod{26} = 3 \pmod{26} = 3$      |
| ↓  | ↓ | ↓  | ↓  | $f(17) = 7 \cdot 17 + 3 \pmod{26} = 122 \pmod{26} = 18$ |
| 4  | 3 | 18 | 21 | $f(10) = 7 \cdot 10 + 3 \pmod{26} = 73 \pmod{26} = 21$  |
| E  | D | S  | V  |   |

## 3. How to decrypt the message encrypted by using affine cipher $f(p) = ap + b \pmod{26}$ ?

\*  $\gcd(a, 26) = 1$ . Let  $p$  be that message,  $c$  be the cipher letter

$$f(p) \equiv a \cdot p + b \pmod{26} = c$$

encryption knowing  $p \rightarrow$  finding  $c$

decryption knowing  $c \rightarrow$  finding  $p$

That is, solving  $p$  in the equation

$$(ap + b) \pmod{26} = c \pmod{26}$$

$$ap \pmod{26} = c - b \pmod{26}$$

we can find inverse of  $a$  modulo 26:  $\bar{a}$  ( $a\bar{a} \equiv 1 \pmod{26}$ )  
( $\gcd(a, 26) = 1$  guarantee the existence of  $\bar{a}$ )

$$\bar{a}ap \pmod{26} = \bar{a}(c - b) \pmod{26}$$

$$p \pmod{26} = \bar{a}(c - b) \pmod{26}$$

$\Rightarrow$  decryption function is  $g(c) = \bar{a}(c - b) \pmod{26}$

4. (a) Find the decryption function for  $f(p) = 7p + 3 \pmod{26}$ . (b) Decrypt the ciphertext message "EDSV" which is encrypted by using  $f(p) = 7p + 3 \pmod{26}$ .

(a) To find the inverse of 7 modulo 26.

$$\begin{aligned} \gcd(7, 26) &= \gcd(7, 26 \bmod 7) \rightarrow 26 = 7 \times 3 + 5 \quad (1) \\ &= \gcd(7, 5) = \gcd(7 \bmod 5, 5) \rightarrow 7 = 5 \times 1 + 2 \quad (2) \\ &= \gcd(2, 5) = \gcd(2, 5 \bmod 2) \rightarrow 5 = 2 \times 2 + 1 \quad (3) \\ &= \gcd(2, 1) = 1 \end{aligned}$$

(2):  $2 = 7 - 5 \times 1$

(3):  $1 = 5 - 2 \times 2 = 5 - 2(7 - 5 \times 1) = 5 - 2 \times 7 + 5 \times 2 = 5 \times 3 - 2 \times 7$

(1):  $5 = 26 - 7 \times 3$

$$\Rightarrow 1 = -11 \times 7 + 26 \times 3$$

$\Rightarrow$  decryption function:

$$g(c) \equiv -11(c-3) \pmod{26} = 15(c-3) \pmod{26}$$

b)

|                                       |                                      |   |   |
|---------------------------------------|--------------------------------------|---|---|
| E                                     | D                                    | S                                       | V                                       |
| 4                                     | 3                                    | 18                                      | 21                                      |
| $\downarrow$                          | $\downarrow$                         | $\downarrow$                            | $\downarrow$                            |
| $g(4) \equiv -11(4-3) \pmod{26} = 15$ | $g(3) \equiv -11(3-3) \pmod{26} = 0$ | $g(18) \equiv -11(18-3) \pmod{26} = 17$ | $g(21) \equiv -11(21-3) \pmod{26} = 10$ |
| P                                     | A                                    | R                                       | K                                       |

$$\begin{aligned} g(4) &\equiv -11(4-3) \pmod{26} = -11 \pmod{26} = 15 \\ g(3) &\equiv -11(3-3) \pmod{26} = 0 \pmod{26} = 0 \\ g(18) &\equiv -11(18-3) \pmod{26} = -165 \pmod{26} = 17 \\ g(21) &\equiv -11(21-3) \pmod{26} = -198 \pmod{26} = 10 \end{aligned}$$

## 5. Public Key Cryptography.

Private key cryptosystems: shift cipher, affine cipher, etc.

Anyone who knows the key can both encrypt and decrypt message which make this cryptosystem simple but extremely vulnerable to cryptanalysis.

Public key cryptosystems: the RSA (Rivest, Shamir, Adleman) system, etc.

One who knows how to encrypt the message does not help decrypt message. Everyone can have a publicly known encryption key but the decryption keys are kept secret.

## 6. RSA Encryption.

Let  $m$  be the message represented in the format of strings and  $(n, e)$  be the public key such that  $n = pq$  is a product of two primes numbers and  $\gcd(e, (p-1)(q-1)) = 1$ . Then we have the encryption function

$$f(m) = m^e \pmod{n}.$$