

# MAT2440, Classwork34, Spring2025

ID: \_\_\_\_\_ Name: \_\_\_\_\_

1. Input (size  $n$ )  $\longrightarrow$  Find the number of operations  $f(n)$   $\longrightarrow$   $f(n) = O(\log n)$   
 measurement of time complexity  $\uparrow$
2. Commonly Used Terminology for the Complexity of Algorithms.

Complexity	Terminology
$O(1)$	<u>constant</u> complexity
$O(\log n)$	<u>Logarithmic</u> complexity
$O(n)$	<u>Linear</u> complexity
$O(n \cdot \log n)$	<u>Linearithmic</u> complexity
$O(n^p), p > 0$	<u>Polynomial</u> complexity
$O(b^n), b > 1$	<u>Exponential</u> complexity
$O(n!)$	<u>Factorial</u> complexity

3.

TABLE 2 The Computer Time Used by Algorithms.						
Problem Size	Bit Operations Used					
$n$	$\log n$	$n$	$n \log n$	$n^2$	$2^n$	$n!$
10	$3 \times 10^{-11}$ s	$10^{-10}$ s	$3 \times 10^{-10}$ s	$10^{-9}$ s	$10^{-8}$ s	$3 \times 10^{-7}$ s
$10^2$	$7 \times 10^{-11}$ s	$10^{-9}$ s	$7 \times 10^{-9}$ s	$10^{-7}$ s	$4 \times 10^{11}$ yr	*
$10^3$	$1.0 \times 10^{-10}$ s	$10^{-8}$ s	$1 \times 10^{-7}$ s	$10^{-5}$ s	*	*
$10^4$	$1.3 \times 10^{-10}$ s	$10^{-7}$ s	$1 \times 10^{-6}$ s	$10^{-3}$ s	*	*
$10^5$	$1.7 \times 10^{-10}$ s	$10^{-6}$ s	$2 \times 10^{-5}$ s	0.1 s	*	*
$10^6$	$2 \times 10^{-10}$ s	$10^{-5}$ s	$2 \times 10^{-4}$ s	0.17 min	*	*

Time needed less  $\longrightarrow$  more  
 \* means time of more than  $10^{100}$  years

4. The Worst-Case Complexity.

The worst-case performance of an algorithm means the largest number of operations needed to solve the given problem using this algorithm on input of specified size. This worst-case analysis tells us how many operations an algorithm requires to guarantee that it will produce a solution.

5. The time complexity analysis of the **linear search** algorithm:

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procedure linear_search( $x$ : integer,  $a_1, a_2, \dots, a_n$ : distinct integers)
 $n$  := the length of  $\{a_i\}$ 
 $i$  := 1
while ( $i \leq n$  and  $x \neq a_i$ )
     $i$  :=  $i + 1$ 
if  $i \leq n$  then  $location$  :=  $i$ 
    else  $location$  := 0
return  $location$  {  $location$  is the subscript of the term that equals  $x$ , or 0 if  $x$  is not found.}
    
```

(1) The number of operations considering the **worst-case**:  $x$  is not in the list

<pre> <b>while</b> (<math>i \leq n</math> and <math>x \neq a_i</math>)     <math>i</math> := <math>i + 1</math> <b>if</b> <math>i \leq n</math> <b>then</b> <math>location</math> := <math>i</math>     <b>else</b> <math>location</math> := 0                 </pre>	$\leftarrow$ <u>2</u> comparison $\leftarrow$ <u>1</u> addition $\leftarrow$ <u>1</u> comparison	$\left. \begin{array}{l} \text{ } \end{array} \right\} n^{\text{times}} \text{ } \left. \begin{array}{l} \text{ } \end{array} \right) 3n$
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$i = 1 \text{ to } n$

$) + 1$

(2) The time complexity is  $f(n) = \underline{3n+1} = O(\underline{n})$  which is linear complexity.

(3) What is the best-case complexity of this algorithm?

Sol: best case: " $x = a_1$ "

then the operations are: while loop  $1 \leq n$  and  $x = a_1$  ② comp.  
 if loop  $1 \leq n$  ① comp.

$f(n) = 3 = O(1) \leftarrow$  constant complexity. 3 operations